
Virtual Circuit, Wireless, Spanning Tree

30 November 2025
Lecture 6

Some Slides Credits: Steve Zdancewic (UPenn), Kurose and Ross

Topics for Today

- Virtual Circuit Routing
- 802.11 Wireless
- Bridges and Spanning Tree Algorithm

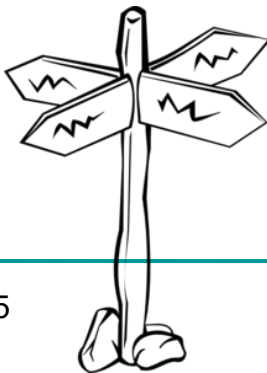
Routing: Drive to Afula

Option 1:

The road authority puts a sign at each intersection to show which way to go

If there are 200 potential destinations, there are 200 arrows at each intersection

At each intersection I must search through 200 arrows to find the one for Afula.

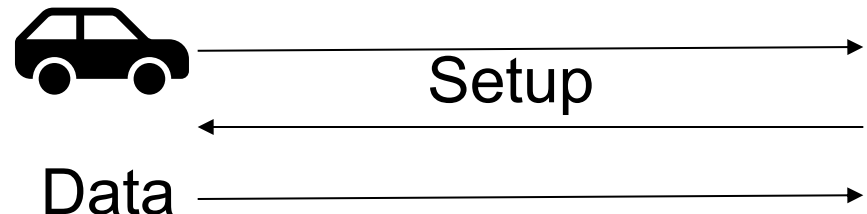


Option 2:

First send someone who knows the way. He puts a sign at each intersection with the name “Michael” and the direction to travel.

There only have to be as many signs at each intersection as there are travelers.

Route setup costs one round trip



Datagram approach

Every packet contains a complete destination address

- Enough information so that any switch can decide where the packet should go.
-

Features of datagram approach

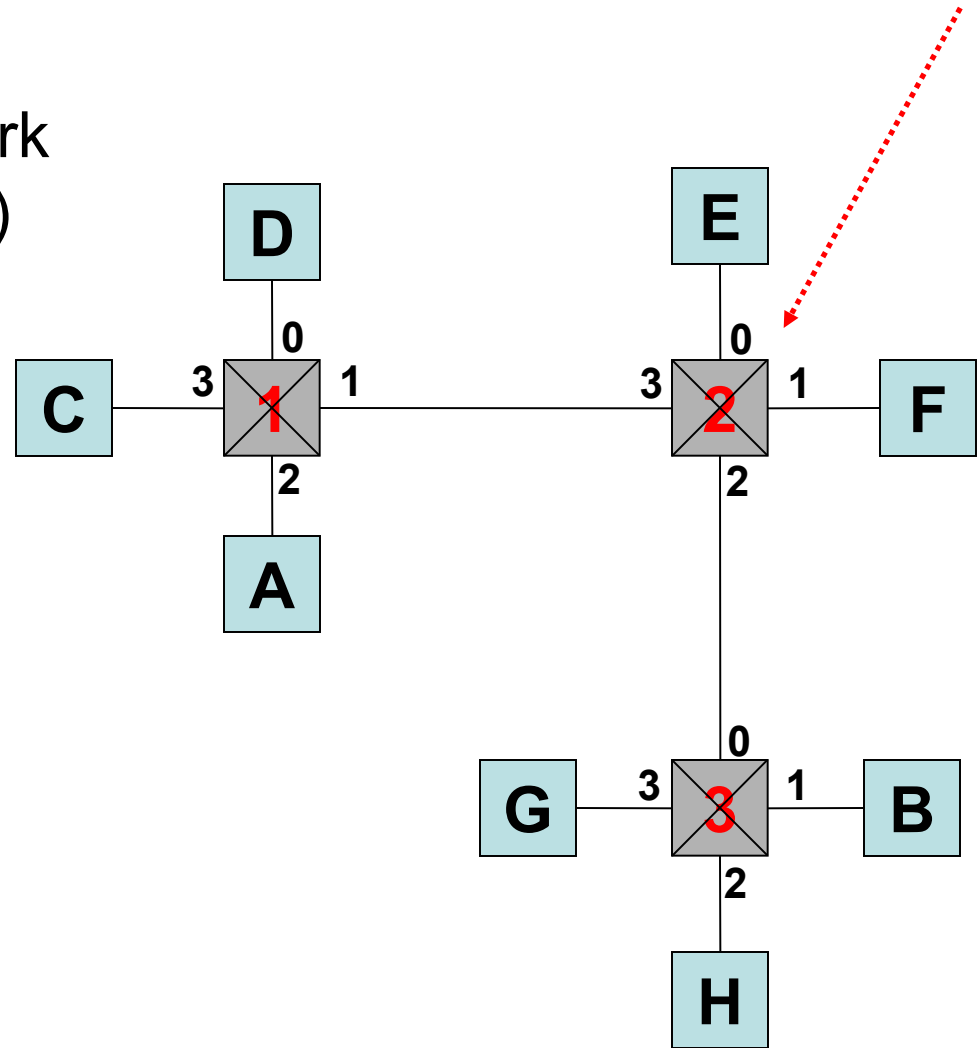
- Packets can be sent *anywhere at any time*
- Sender *doesn't know* if network can deliver the packet (or if destination host is *available*)
- Each packet is forwarded *independently* (two packets may take *different routes*)
- Possible to *route around* switch or link failures

Forwarding Tables

- Provide route information
- Easy to determine if network in known (and unchanging)

Forwarding table
for switch 2

Dest.	Port
A	3
B	2
C	3
D	3
E	0
F	1
G	2
H	2



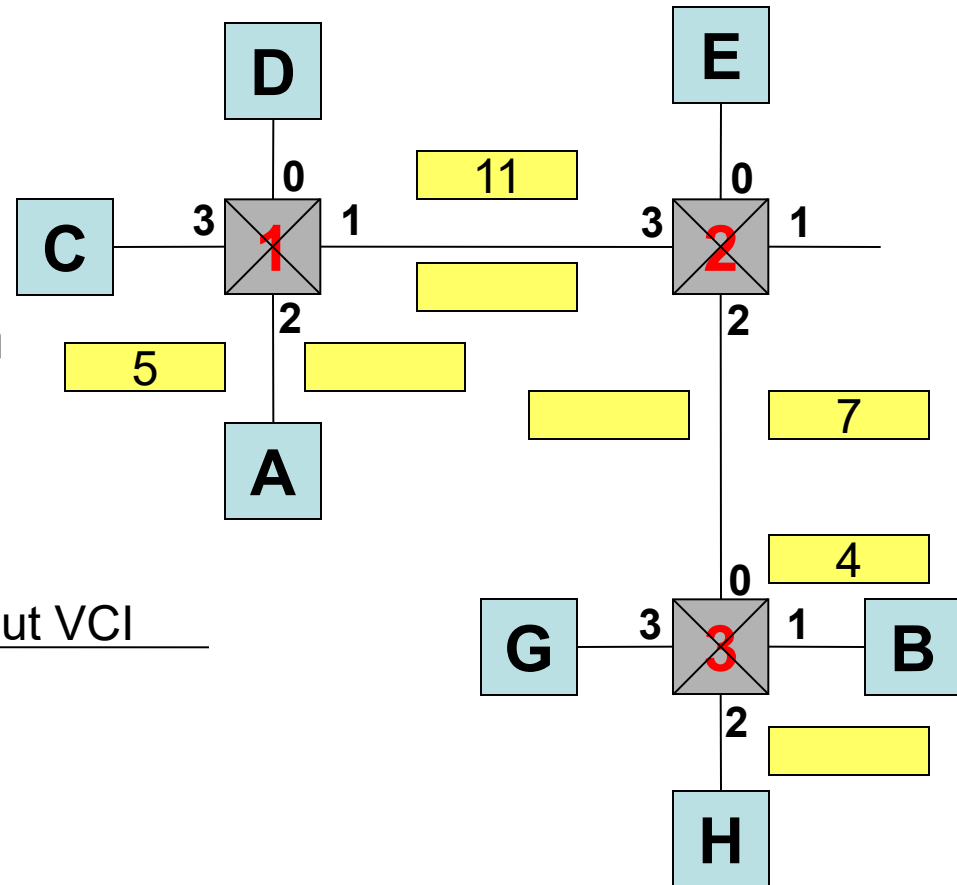
Virtual circuit approach

- Set up the connection **before data transfer**
 - Allocate resources on circuits
 - Set up forwarding tables
- Benefits of virtual circuit approach
 - Performance: per-packet switching cost is low
 - Reliability: predictable latency and throughput
- Drawbacks
 - Setup time is long
 - At least one RTT – why?
 - Fault tolerance
 - What if the circuit fails during the transmission?

Virtual Circuit Switching

- VCI = Virtual Circuit Identifier
- Incoming port + VCI uniquely identify virtual circuit
- Setup phase constructs circuit table entries at each switch

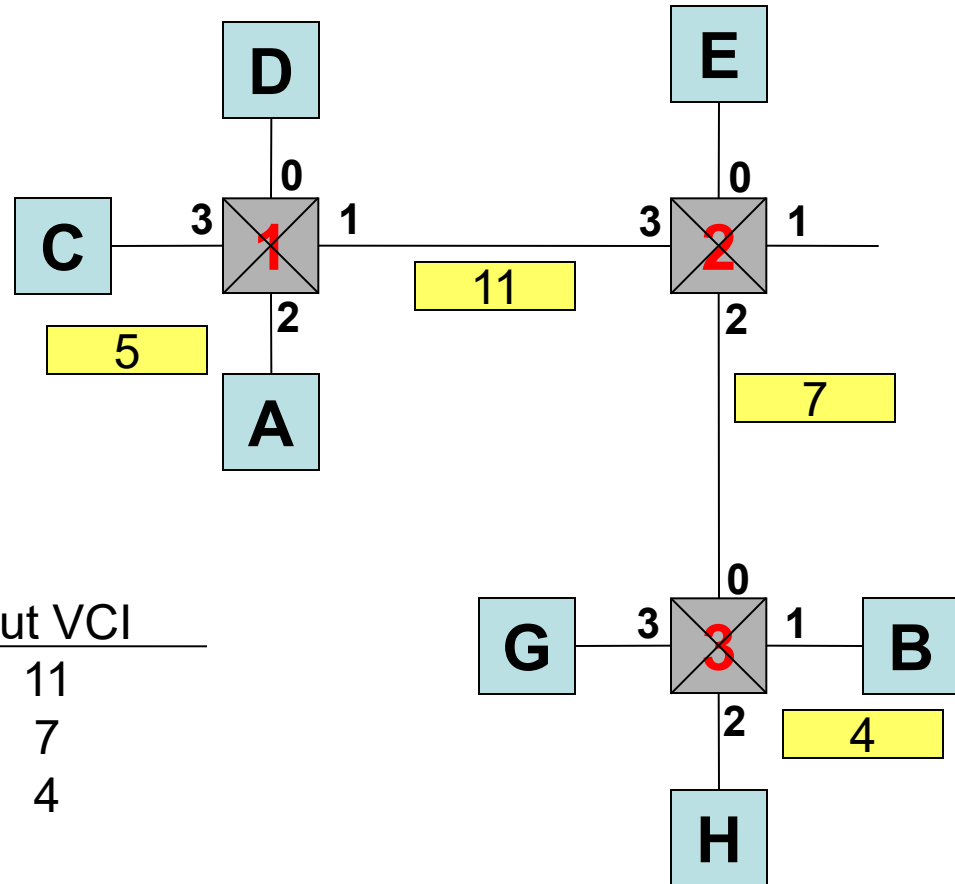
A wants to send to B



Switch	In Port	In VCI	Out port	Out VCI
1				
2				
3				

Virtual Circuit Switching

- VCI = Virtual Circuit Identifier
- Incoming port + VCI uniquely identify virtual circuit
- Setup phase constructs circuit table entries at each switch



Switch	In Port	In VCI	Out port	Out VCI
1	2	5	1	11
2	3	11	2	7
3	0	7	1	4

Datagram versus Virtual Circuit

Datagram

Advantages:

1. Messages have no setup cost.
2. Routing table size depends on the number of nodes, not number of conversations.
3. Faster recovery from network failures.

Disadvantages:

1. Networks with many nodes have slow table lookup.
2. Packet 2 takes just as long to route as packet 1.

Virtual Circuit

Advantages:

1. Routing table size depends on number of conversations.
2. Can configure the circuit once and future messages can route very fast.
3. Save space in packet header.

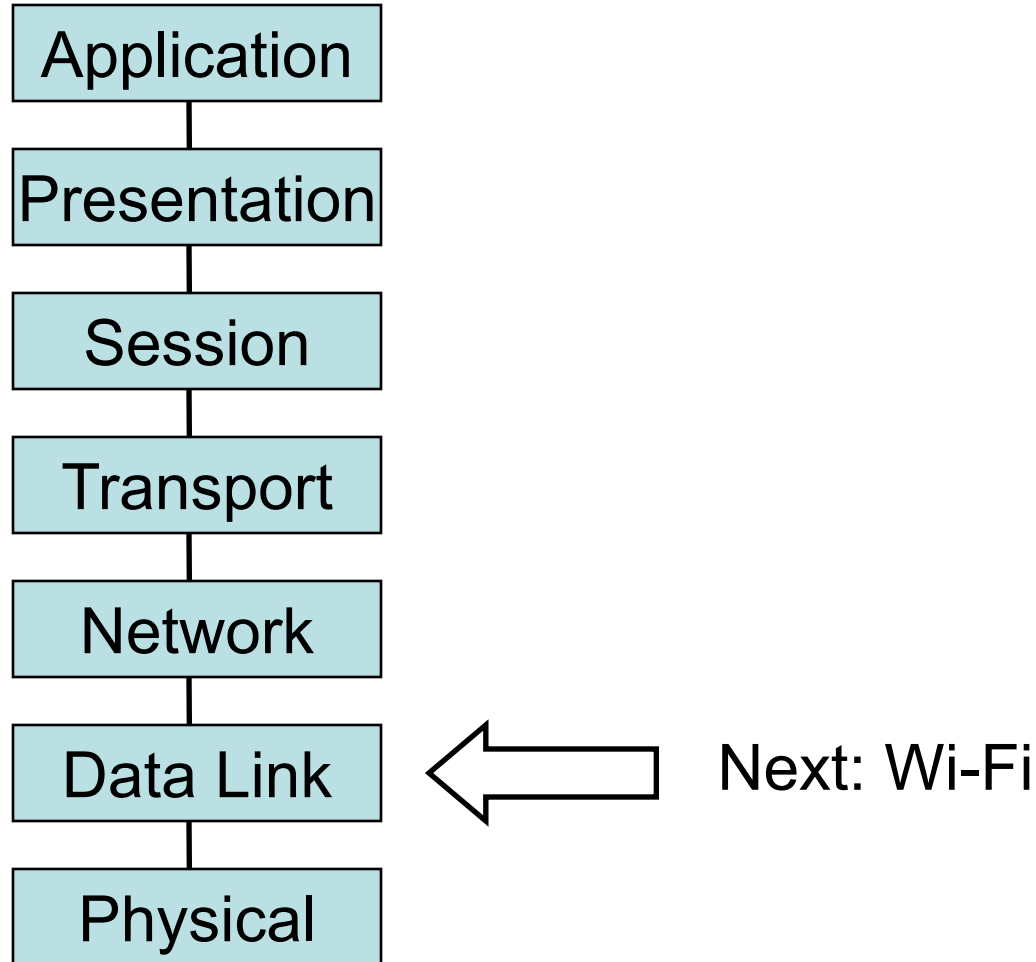
Disadvantages:

1. Dynamic setup is costly (1 round trip).
2. Slower to recover from failures.
3. Can't easily route around problems.
4. Bootstrapping requires existing tables.

So Far

- Virtual Circuit Routing
- 802.11 Wireless
- Bridges and Spanning Tree Algorithm

OSI Reference Model



Wireless (802.11)

Like Ethernet, 802.11 has shared medium

- Need MAC
- Uses exponential backoff

Unlike Ethernet, in 802.11

- No support for collision detection
- Not all senders and receivers are directly connected

Background

Number of **wireless**
(mobile) phone
subscribers now exceeds
the number of **wired**
phone subscribers!

Computer networks

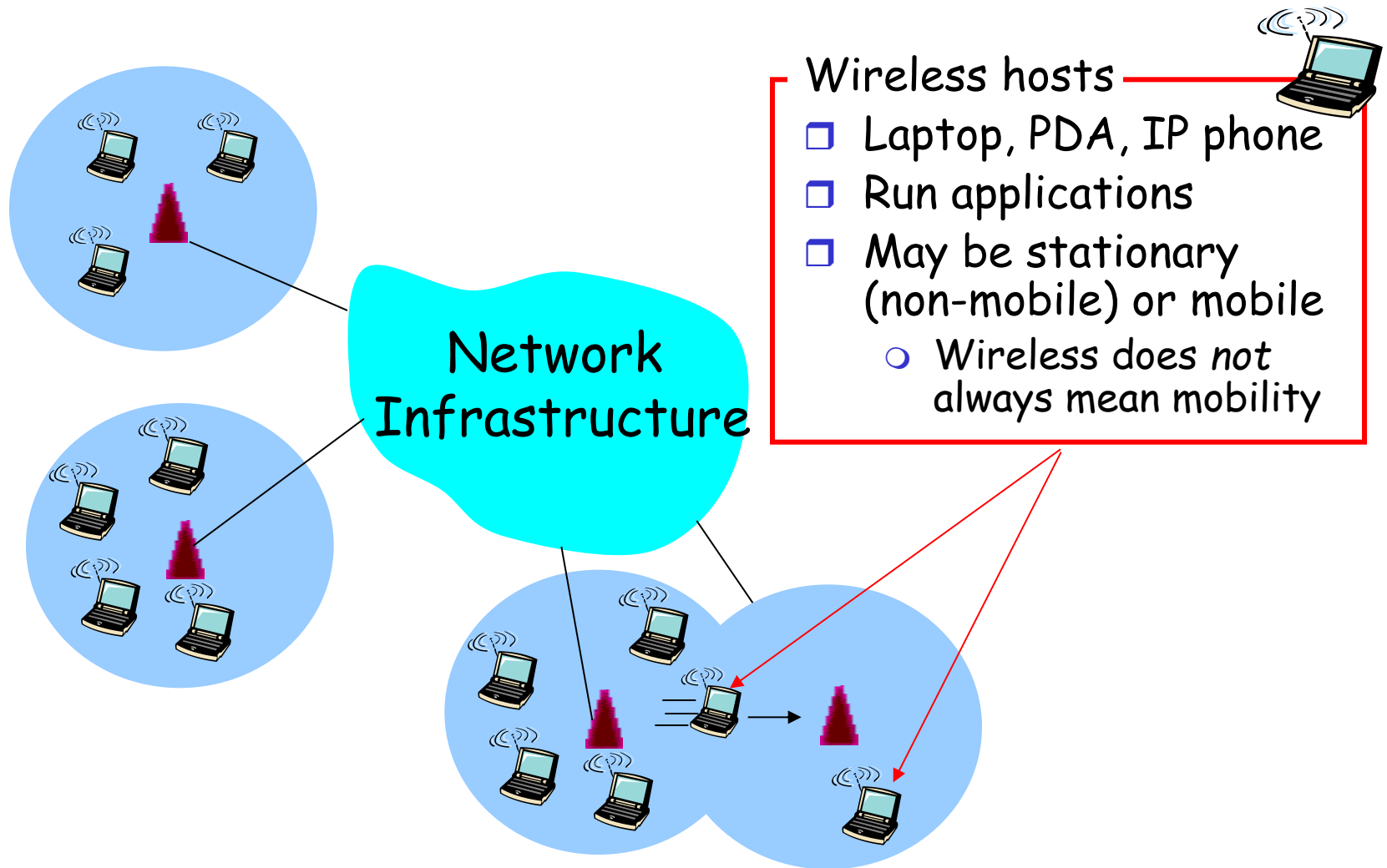
- laptops
- palmtops
- PDAs
- Internet-enabled phones

Promises anytime
untethered Internet
access

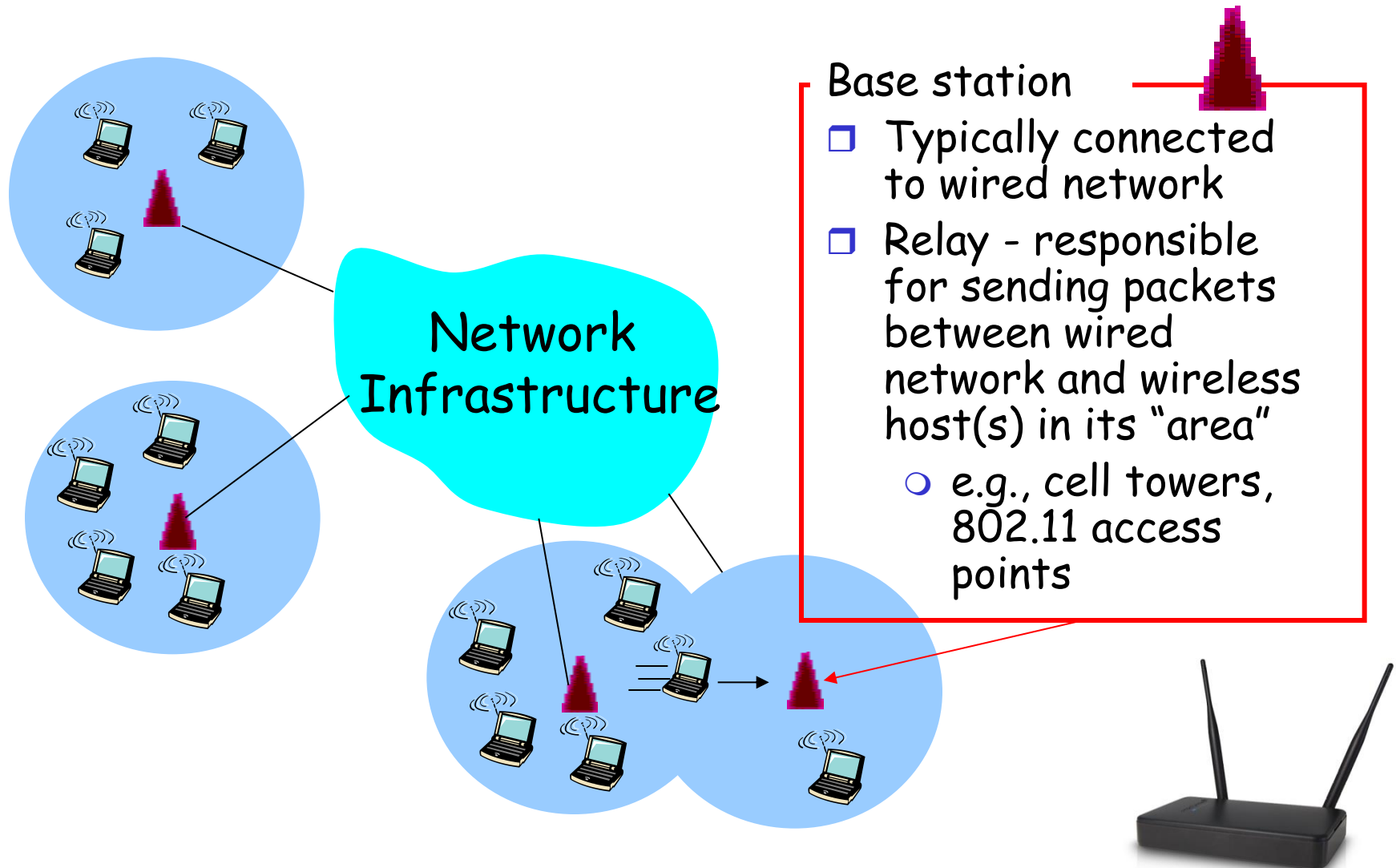
Two important (but different) challenges

- **Wireless**: communication over wireless link
- **Mobility**: handling the mobile user who changes point of attachment to network

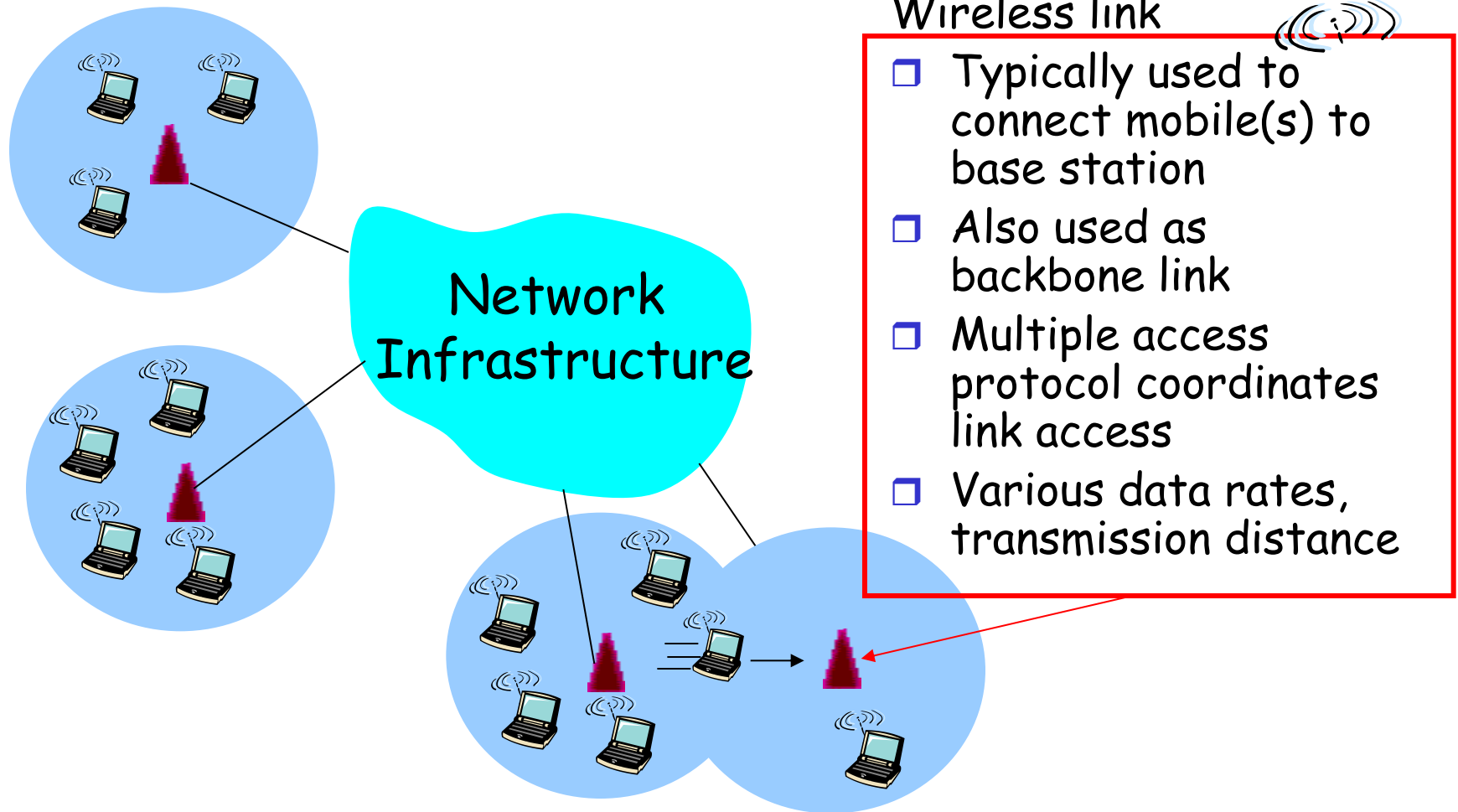
Elements of a Wireless Network



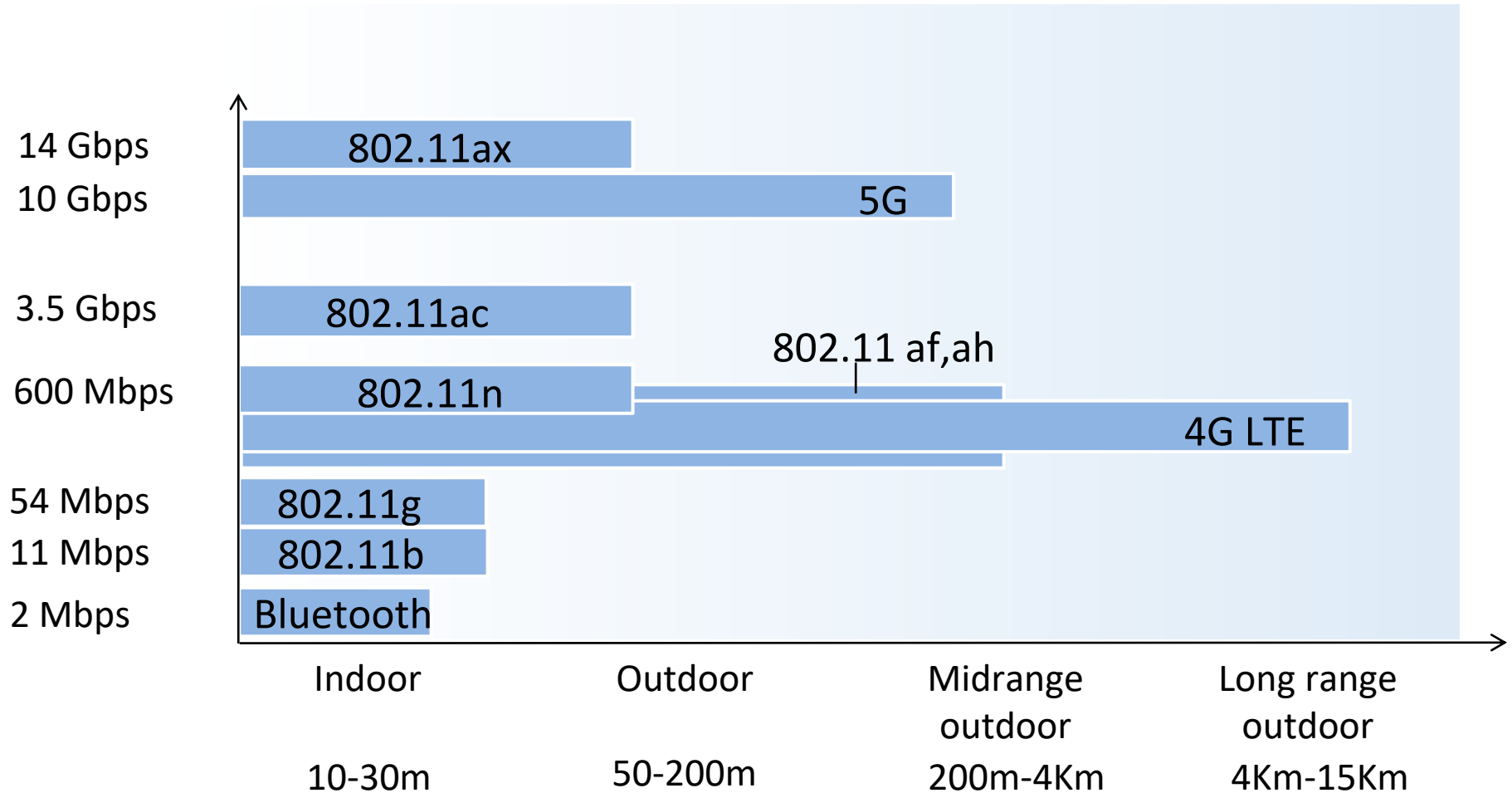
Elements of a Wireless Network



Elements of a Wireless Network



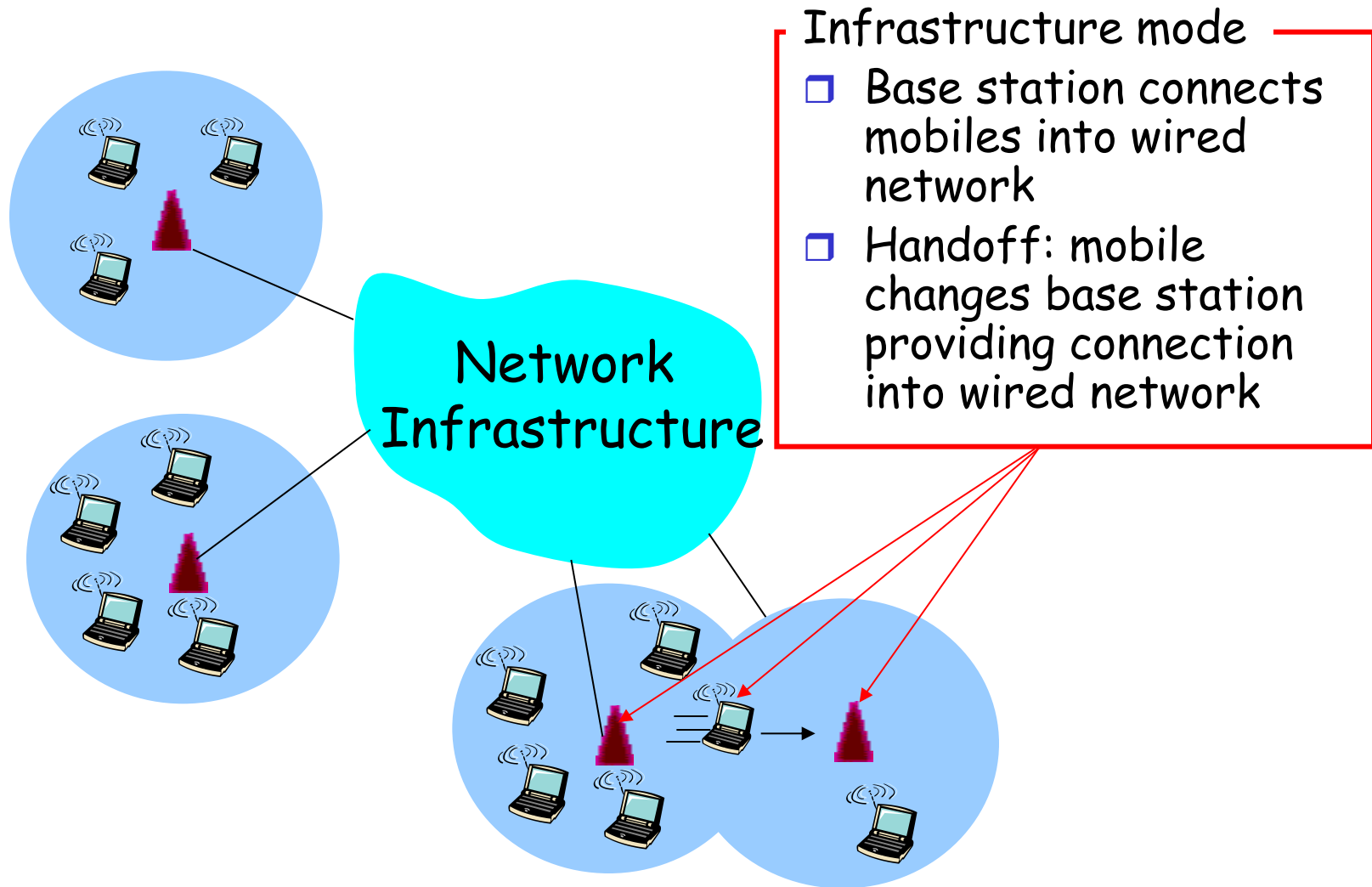
Characteristics of selected wireless link standards



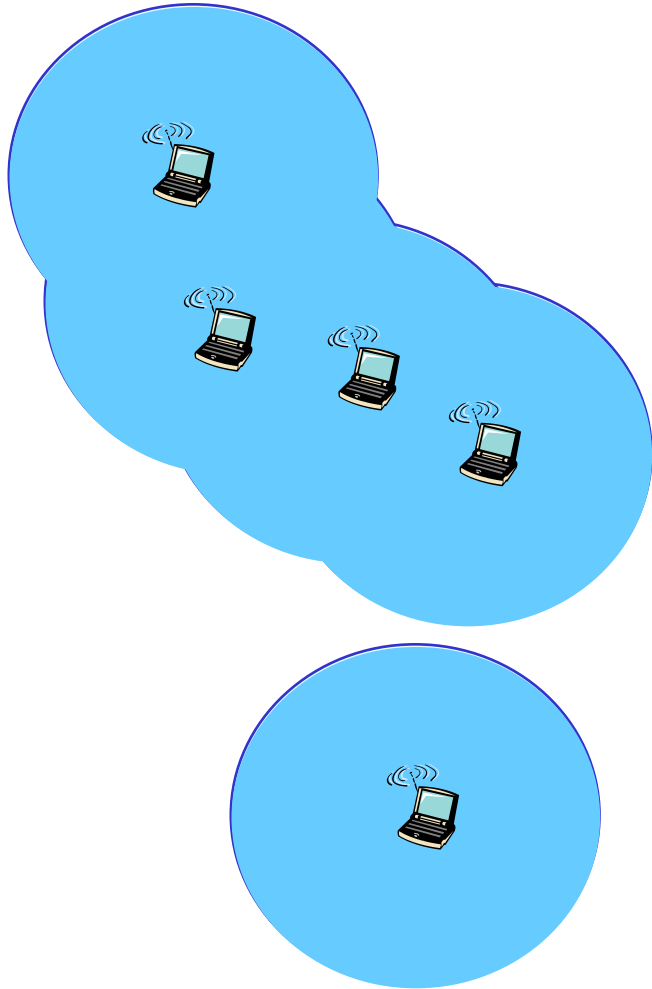
802.11 Wireless LAN Standards

IEEE 802.11 standard	Year	Max data rate	Range	Frequency
802.11b	1999	11 Mbps	30 m	2.4 Ghz
802.11g	2003	54 Mbps	30m	2.4 Ghz
802.11n (WiFi 4)	2009	600	70m	2.4, 5 Ghz
802.11ac (WiFi 5)	2013	3.47 Gpbs	70m	5 Ghz
802.11ax (WiFi 6)	2021	14 Gbps	70m	2.4, 5 Ghz
802.11af	2014	35 – 560 Mbps	1 Km	unused TV bands (54-790 MHz)
802.11ah	2017	347Mbps	1 Km	900 Mhz
802.11be (WiFi 7)	2024	0.4-23,059 Mbps	30-120m	2.4, 5, 6
802.11bn (WiFi 8)	2028 (exp.)	0.4-23,059 Mbps	30-120m	2.4, 5, 6

Elements of a Wireless Network



Elements of a Wireless Network



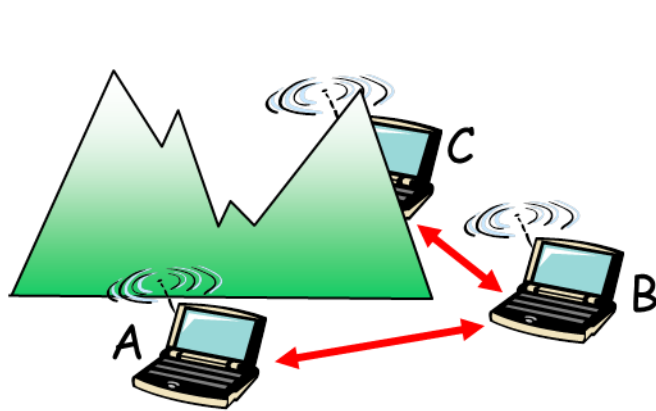
Ad hoc mode

- ☐ No base stations
- ☐ Nodes can only transmit to other nodes within link coverage
- ☐ Nodes organize themselves into a network: route among themselves

Wireless Network Taxonomy

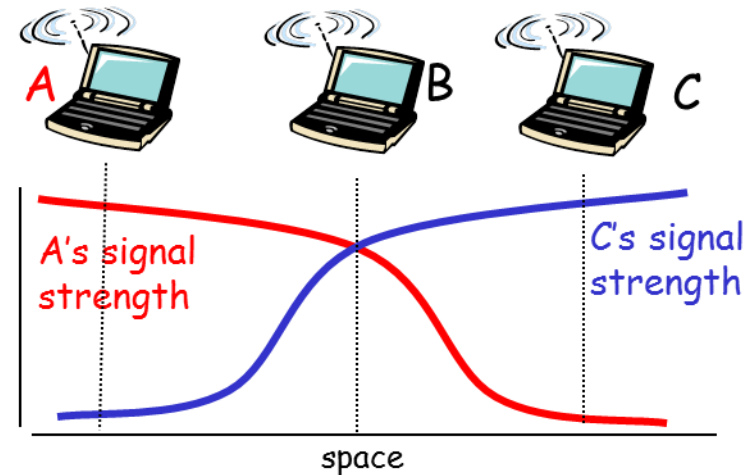
	Single Hop	Multiple Hops
Infrastructure	host connects to base station (WiFi, WiMAX, cellular) which connects to larger Internet	host may have to relay through several wireless nodes to connect to larger Internet: mesh net
No Infrastructure	no base station, no connection to larger Internet (Bluetooth, ad hoc nets)	no base station, no connection to larger Internet. May have to relay to reach other a given wireless node MANET, VANET

Hidden Node Problem



Hidden terminal problem

- ☐ B, A hear each other
 - ☐ B, C hear each other
 - ☐ A, C can not hear each other
- means A, C unaware of their interference at B



Signal attenuation:

- ☐ B, A hear each other
 - ☐ B, C hear each other
 - ☐ A, C can not hear each other
- interfering at B

Exposed Node Problem



Exposed terminal problem

- ☐ B, C hear each other
- ☐ B wants to talk to A
- ☐ C wants to talk to D
- ☐ B needs to know that D can't hear B and is OK to send

Signal attenuation:

- ☐ B, C hear each other
- ☐ A, C don't hear each other
- ☐ D, B don't hear each other

How 802.11 works

Medium is *shared*

Collision domains
are more complex

Method of operation:
CSMA/CA

- *Carrier sensing multiple access, with **collision avoidance***

Augmented media
access control
(MAC) protocol:

- Slot reservation protocol

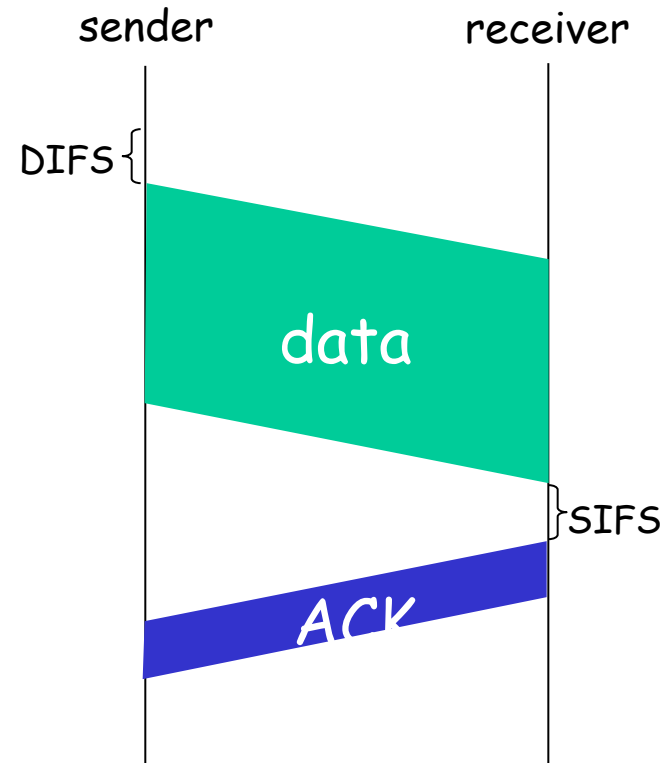
MAC Protocol: CSMA/CA

802.11 sender

- 1 if sense channel idle for **DIFS** then
transmit entire frame (no CD)
- 2 if sense channel busy then
start random backoff time
timer counts down while channel
idle
transmit when timer expires
if no ACK, increase random
backoff interval, repeat 2

802.11 receiver

- if frame received OK
return ACK after **SIFS** (ACK needed
due to hidden terminal problem)



SIFS/DIFS Numbers

Standard	SIFS	Slot Time	DIFS = SIFS+ 2 × Slot Time
IEEE 802.11-1997 (FHSS)	$28\mu s$	$50\mu s$	$128\mu s$
IEEE 802.11-1997 (DSSS)	$10\mu s$	$20\mu s$	$50\mu s$
IEEE 802.11b	$10\mu s$	$20\mu s$	$50\mu s$
IEEE 802.11a	$16\mu s$	$9\mu s$	$34\mu s$
IEEE 802.11g	$10\mu s$	9 or $20\mu s$	28 or $50\mu s$
IEEE 802.11n (2.4 GHz)	$10\mu s$	9 or $20\mu s$	28 or $50\mu s$
IEEE 802.11n (5 GHz)	$16\mu s$	$9\mu s$	$34\mu s$
IEEE 802.11ac	$16\mu s$	$9\mu s$	$34\mu s$

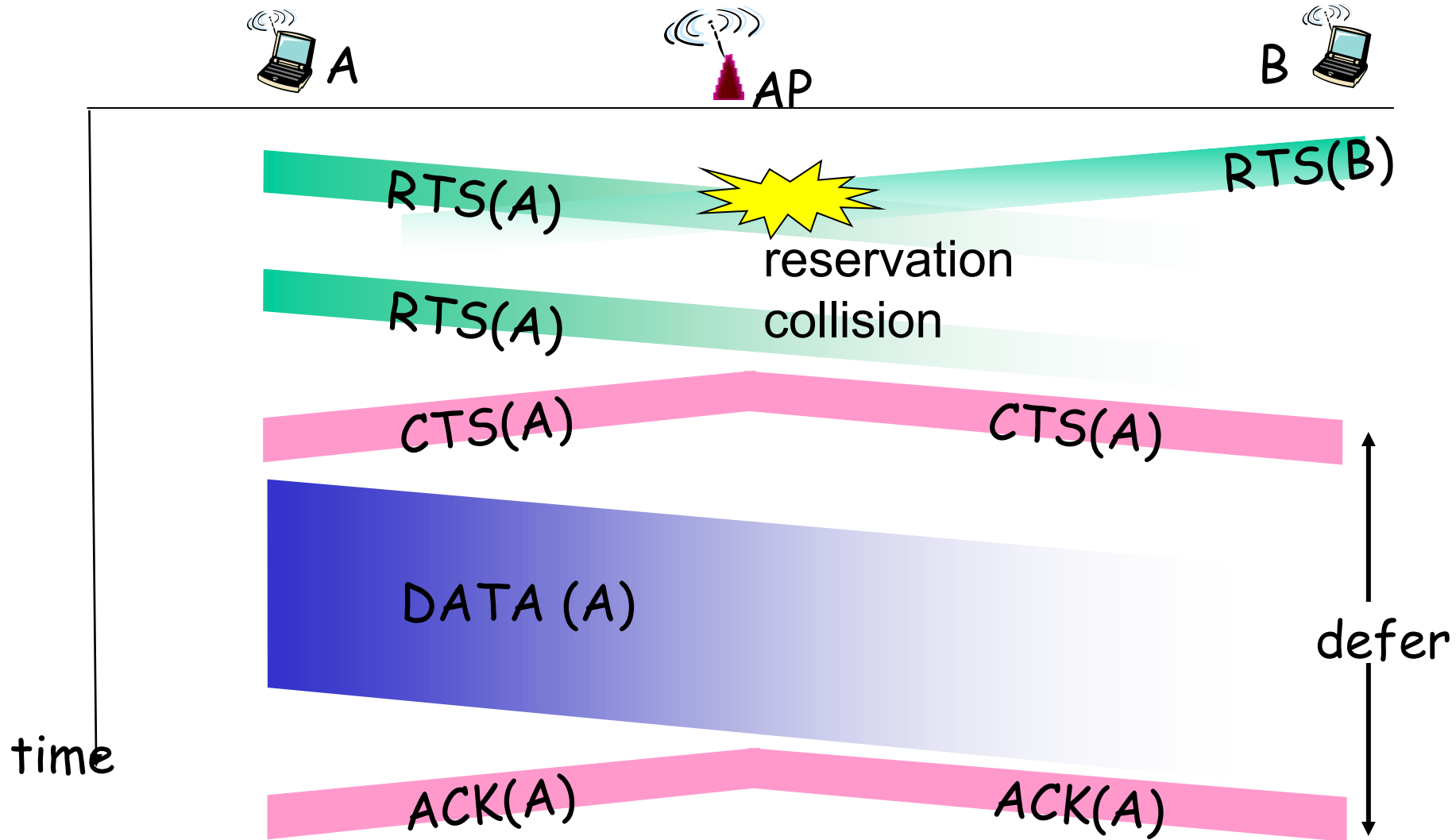
Avoiding Collisions

Idea: Allow sender to “reserve” channel rather than random access of data frames: avoid collisions of long data frames

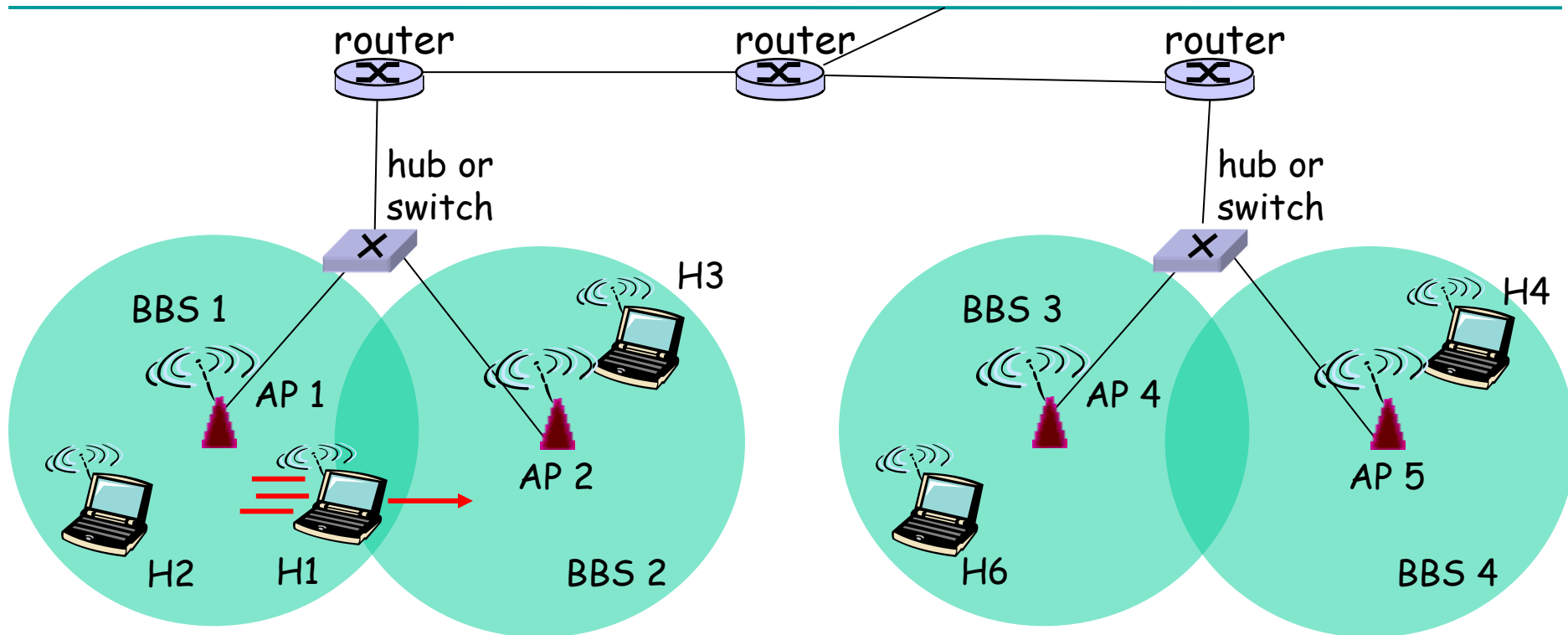
- Sender first transmits *small* request-to-send (RTS) packets to BS using CSMA
 - RTS contains requestor name and length of data to send
 - RTSs may still collide with each other (but they're short)
- BS broadcasts clear-to-send CTS in response to RTS
 - Echoes approved node and the data length to send
- CTS heard by all nodes
 - sender transmits data frame
 - other stations defer transmissions

avoid data frame collisions completely
using small reservation packets!

Collision Avoidance: RTS-CTS

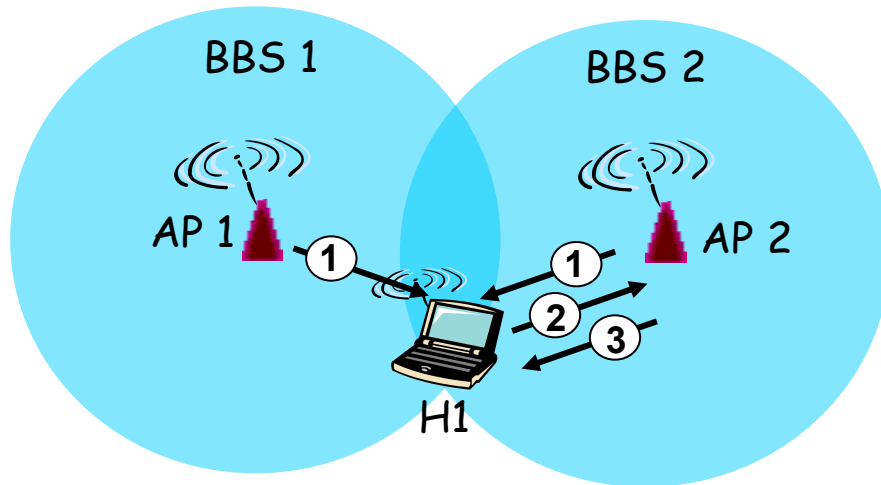


Wireless Access Points



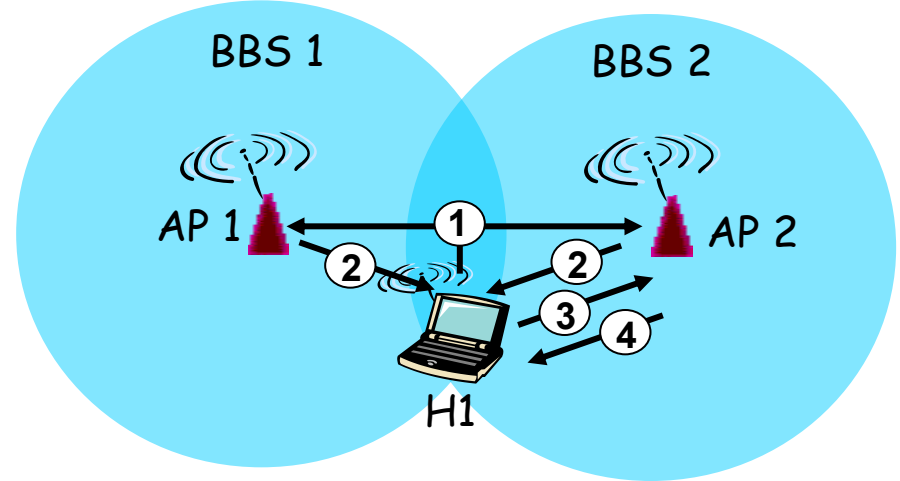
- **Distribution System** – wired network infrastructure connects routers
- **Mobility**: H1 moves right, sees AP2 getting stronger than AP1, requests to associate with AP2

Scanning: Active and Passive



Passive Scanning:

1. Beacon frames sent from APs
2. Association Request frame sent: H1 to selected AP
3. Association Response frame sent: H1 to selected AP



Active Scanning:

1. Probe Request frame broadcast from H1
2. Probes response frame sent from APs
3. Association Request frame sent: H1 to selected AP
4. Association Response frame sent: H1 to selected AP

802.11 Security Issues



Packet Sniffing is *worse*

- No physical connection needed
- Long range (6 blocks)
- Old encryption standards (WEP, WEP2) were bad

Denial of service

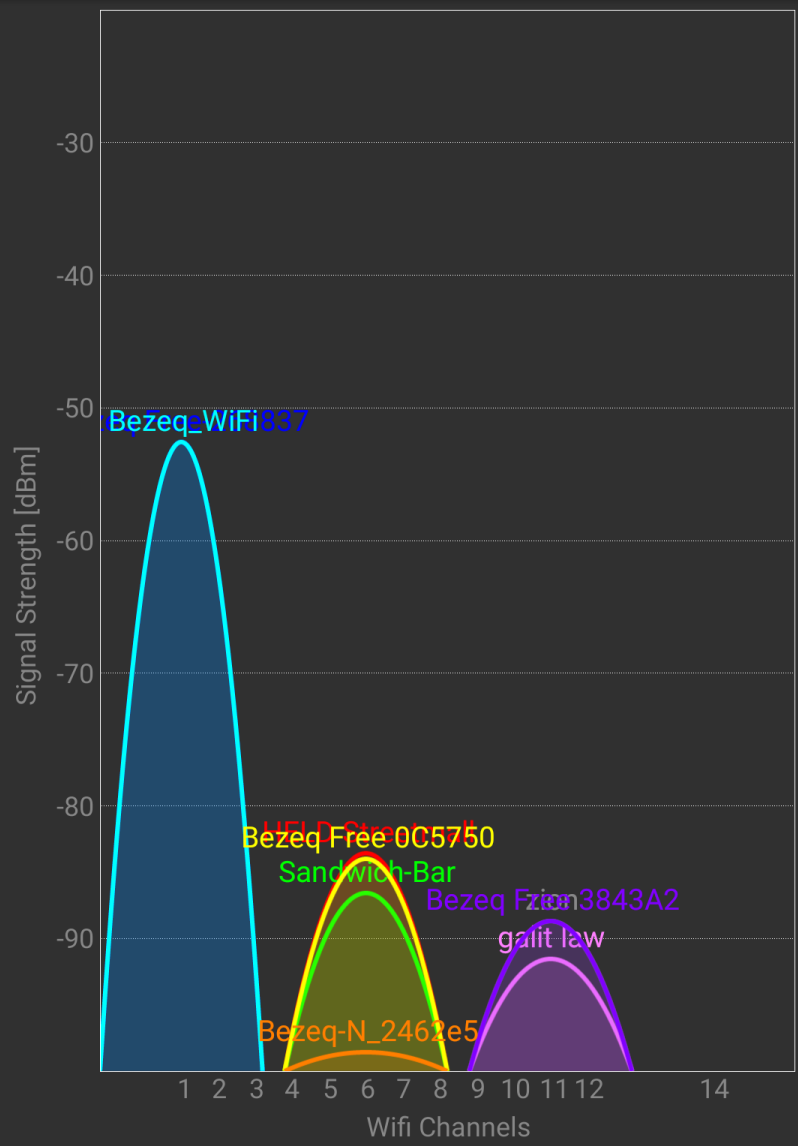
- Association (and Disassociation) Requests are not authenticated

Better: WPA “Wi-Fi Protected Access”

- Introduced in 802.11i
- Uses much stronger cryptology (AES)

More about this in the
SE course
Communication and
Information Security

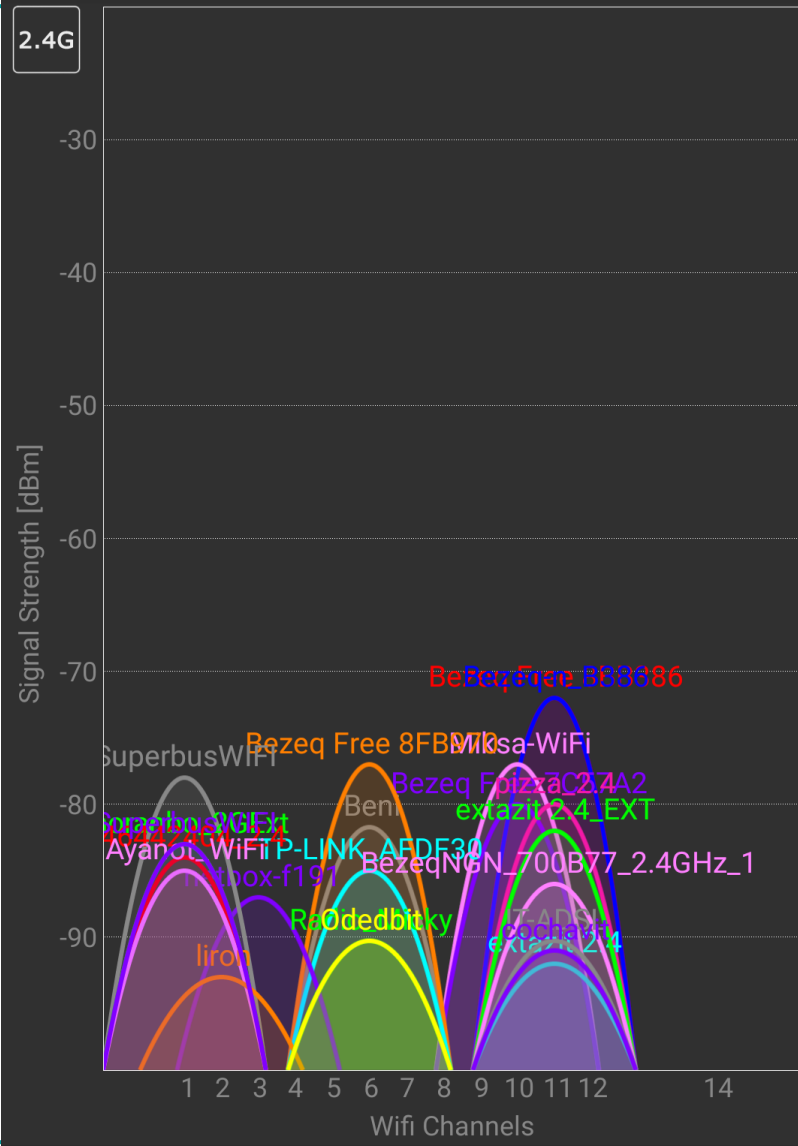
Wifi Analyzer



KS



Wifi Analyzer



Wi-Fi Channels

Two main Wi-Fi frequency zones

- 2.4GHz
- 5GHz

Each zone divided into channels

Hosts and AP communicate over selected channel

- If 2 hosts send on overlapping channels, neither one is understandable

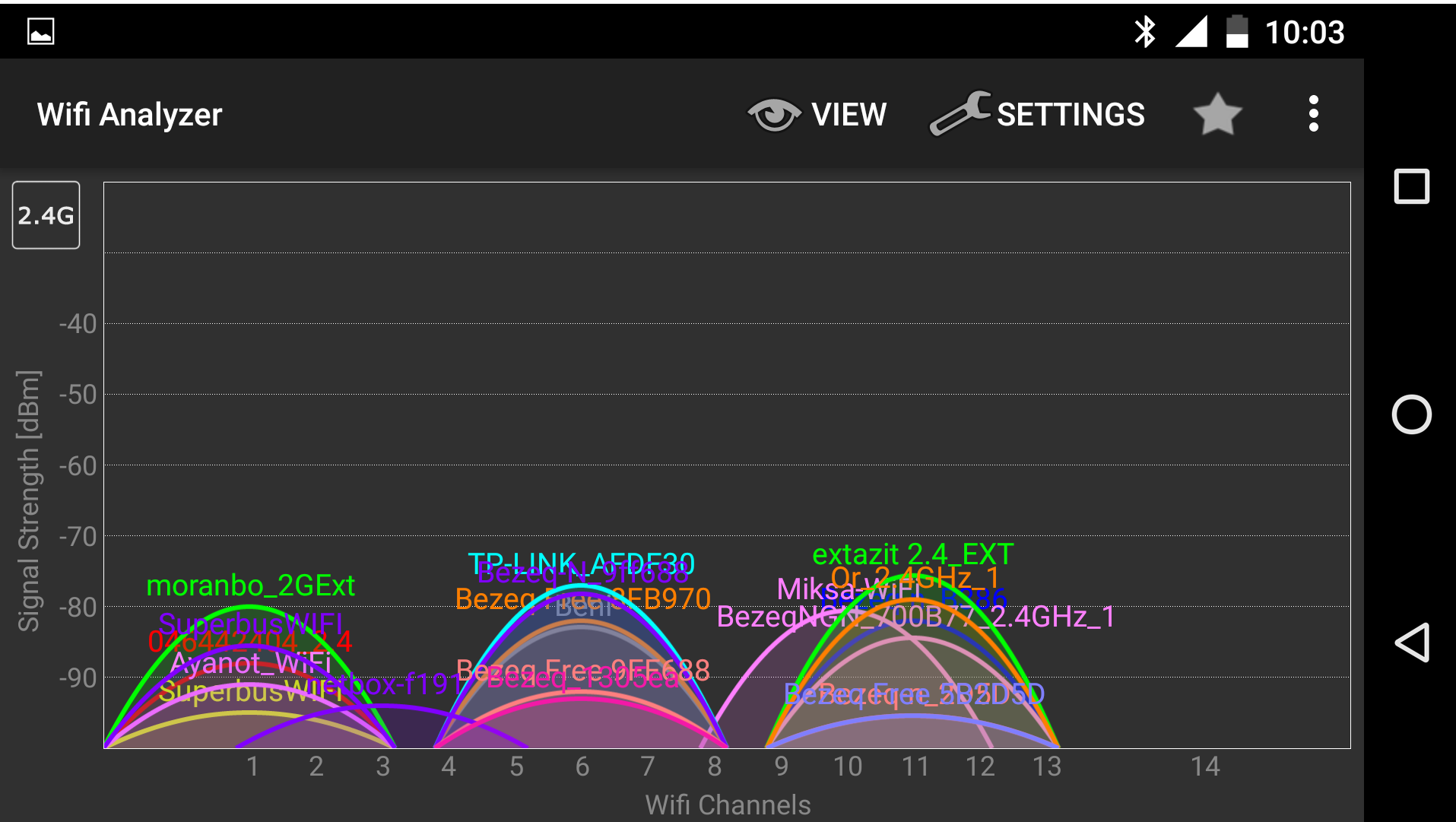
2.4GHz: Lots of overlapping channels

- Can choose multiple non-overlapping ones

5GHz: Less overlap

- Can bond multiple channels for a single message to increase throughput

Wi-Fi 2.4GHz channels



2.4GHz channels

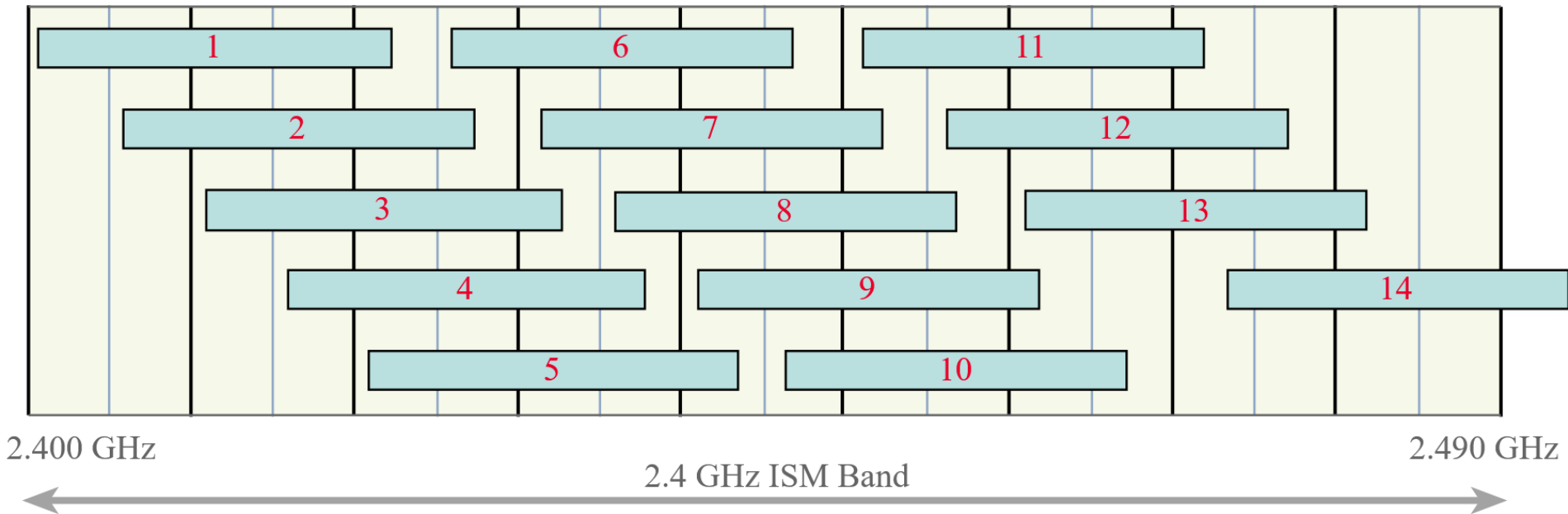
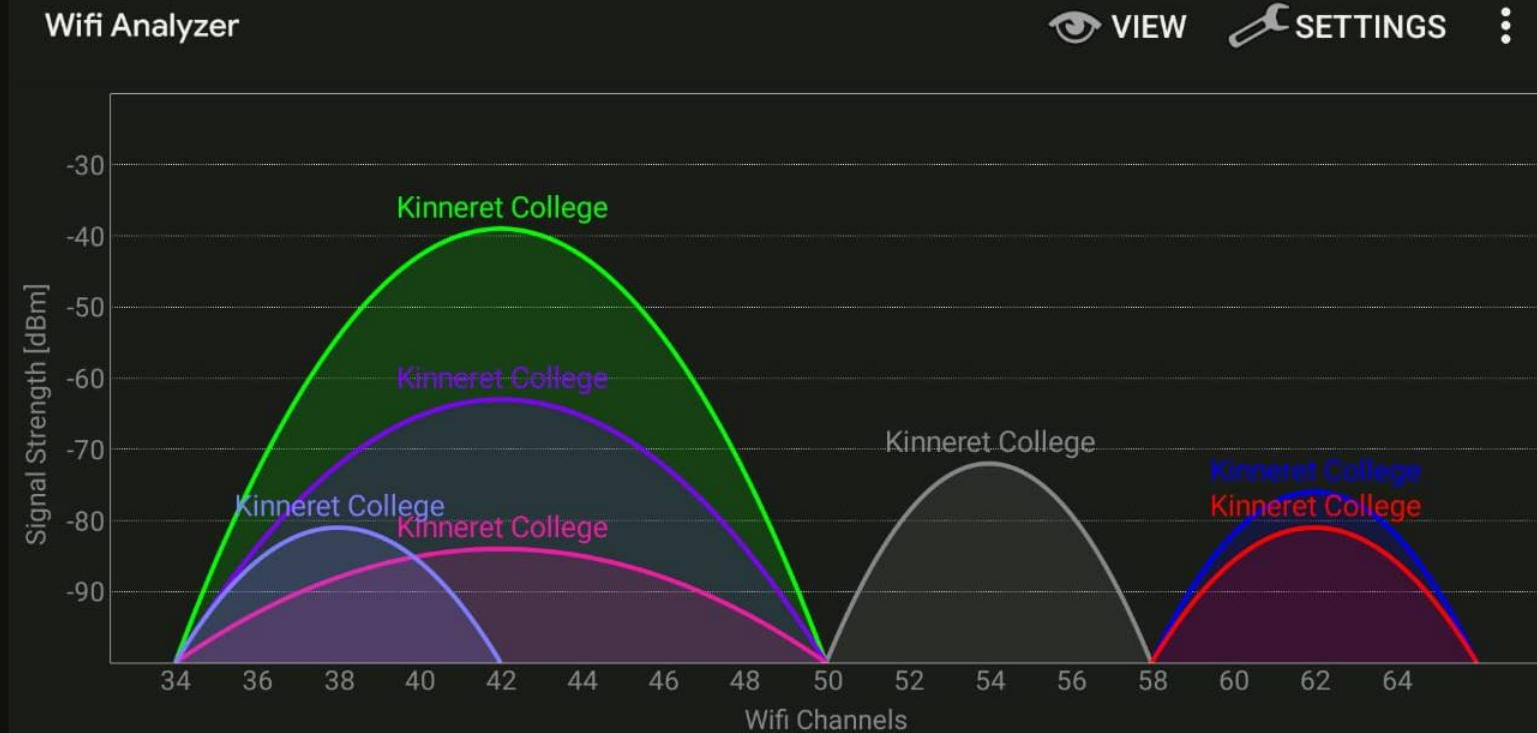


Image source: <https://www.electronics-notes.com/articles/connectivity/wifi-ieee-802-11/channels-frequencies-bands-bandwidth.php>

Wi-Fi 5GHz channels



5GHz channels at 20MHz wide

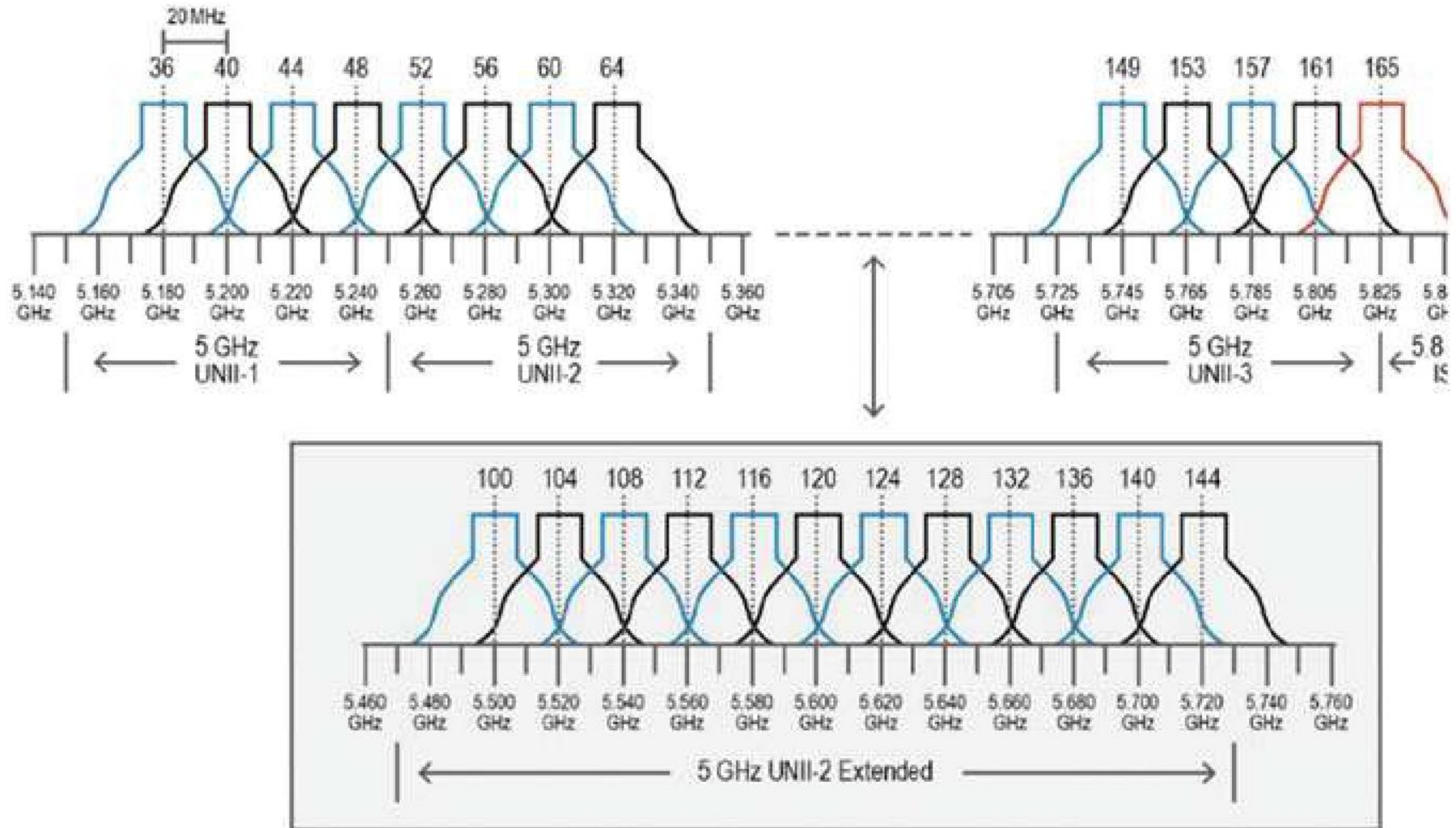
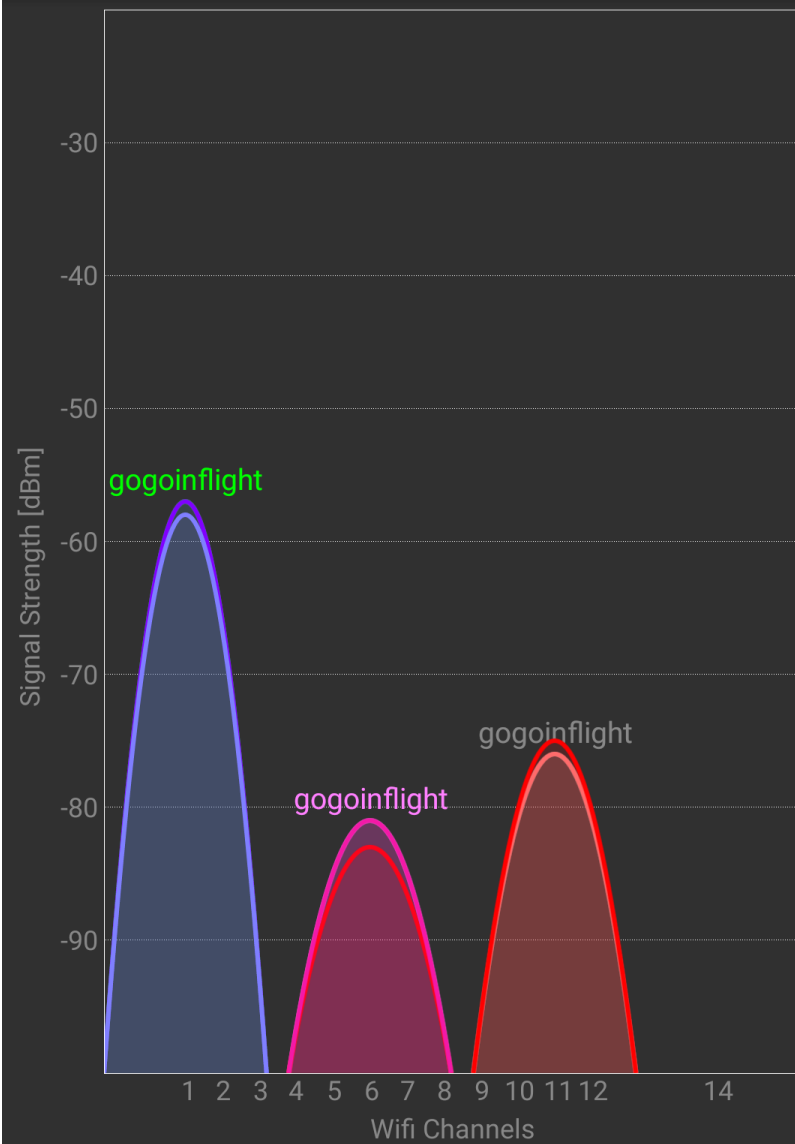


Fig. 3 WLAN channels in 5 GHz band [6]

Lepaja, Salem & Maraj, Arianit & Berzati, Shpat. (2019). WLAN Planning and Performance Evaluation for Commercial Applications: Evolutions in Business Information Processing and Management—Volume 1. 10.1007/978-3-319-94117-2_3.

Wifi Analyzer



Wifi Analyzer

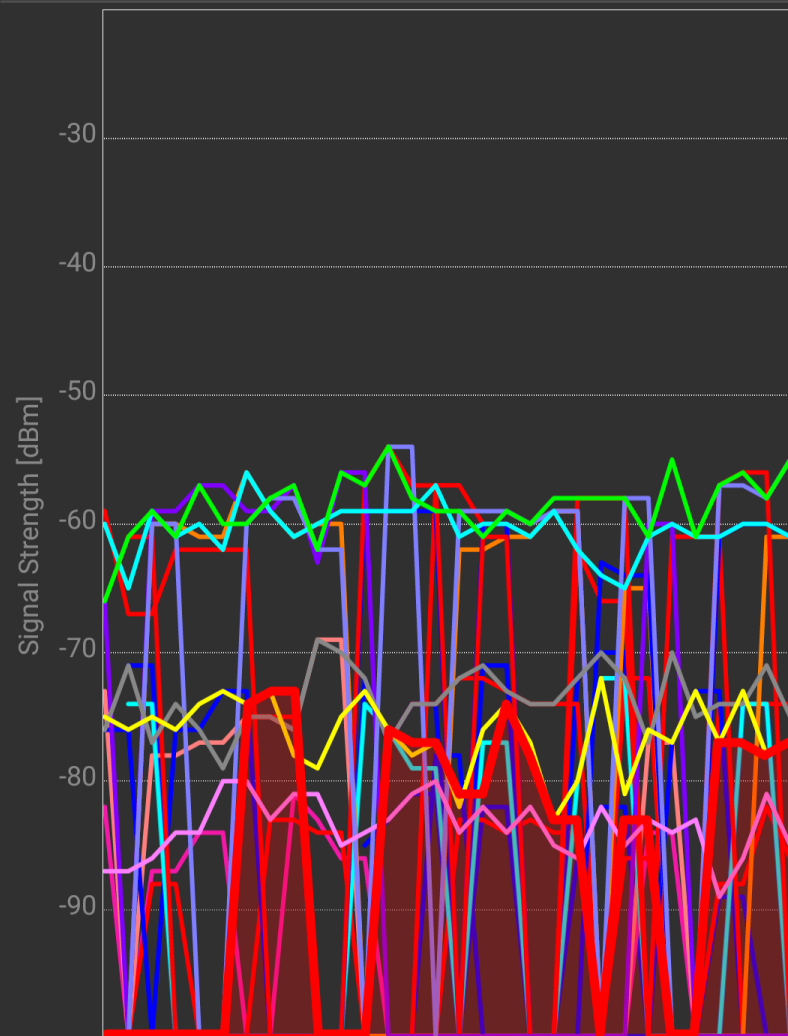


ght

gogoinflight

gogoinflight

gogoinflight



Plane WiFi



Wifi Analyzer

! Not connected!

▶ gogoinflight (...)

CH ... -58 dBm

[ESS]

▶ ? (...)

CH ... -58 dBm

[WPA2-PSK-CCMP][ESS]

▼ ? (...)

CH ... -63 dBm

[WEP][ESS]

? (5c:a4:8a:d8:59:d2)

CH 44 5220 MHz -72 dBm

CISCO SYSTEMS, INC

[WEP][ESS]

? (c0:7b:bc:92:79:92)

CH 36 5180 MHz -63 dBm

CISCO SYSTEMS, INC

[WEP][ESS]

So Far

- Virtual Circuit Routing
- 802.11 Wireless
- Bridges and Spanning Tree Algorithm

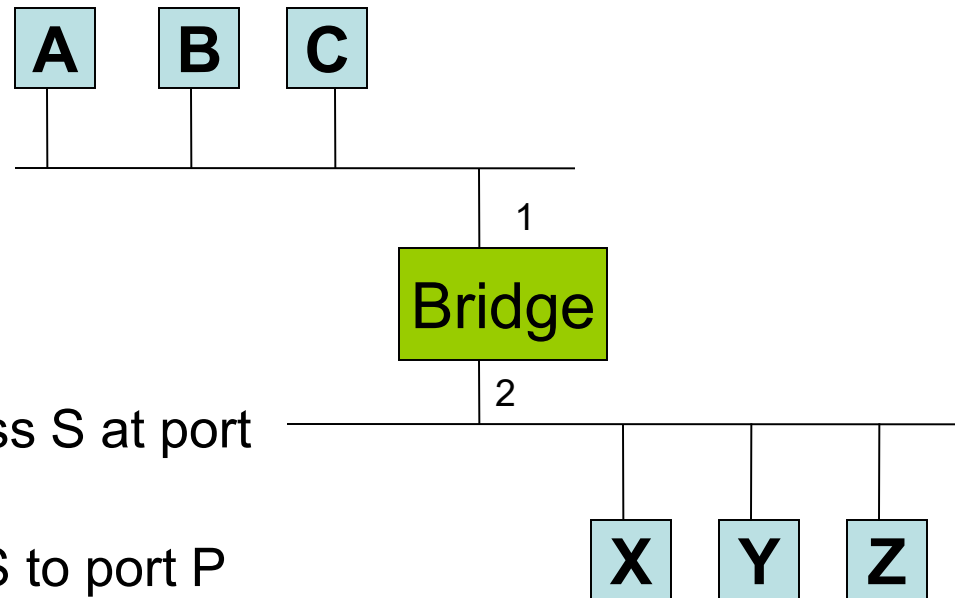
Bridges and LAN Switches

- Bridge accepts LAN frames Bridge on one port, outputs them on another.

- Optimization: only forward appropriate frames

- Learning bridge

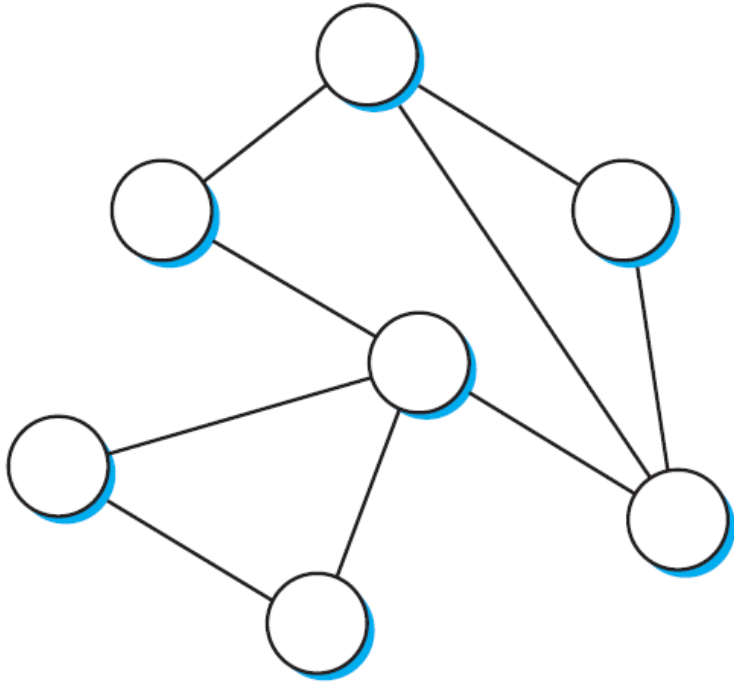
- Watch incoming *source* address S at port number P
- Add entry to forward address S to port P
- If no entry, broadcast to all ports



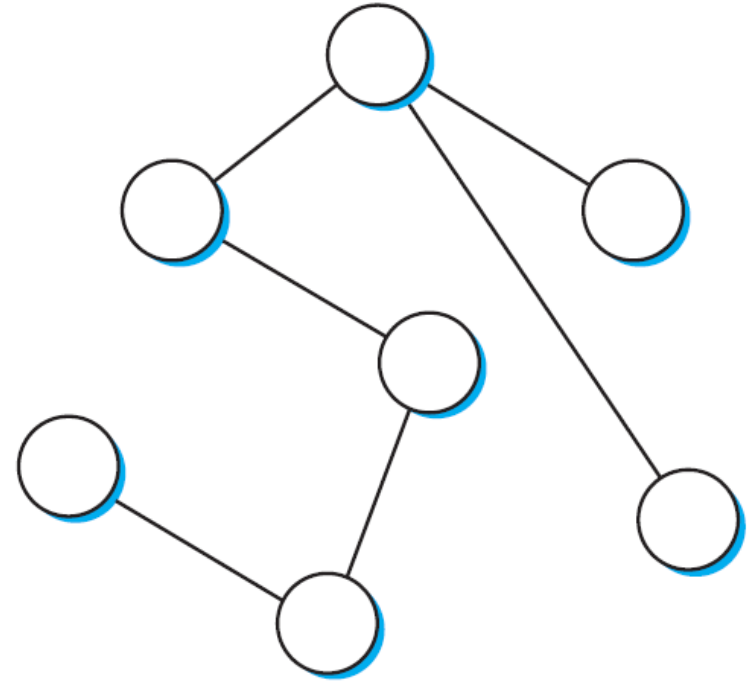
Problem: Cycles (Loops)

- Frame gets rebroadcast forever
- Could avoid by construction, BUT:
 - Hard, especially management
 - Often want redundancy
- Solution:
 - Restrict active ports to a *Spanning Tree*
 - Basic design by Radia Perlman of Digital
 - 802.1 specification of LAN Bridges is based on this algorithm

What is a Spanning Tree?



(a)



(b)

Algorhyme

I think that I shall never see
a graph more lovely than a tree.
A tree whose crucial property
is loop-free connectivity.
A tree that must be sure to span
so packets can reach every LAN.
First, the root must be selected.
By ID, it is elected.
Least-cost paths from root are
traced.
In the tree, these paths are placed.
A mesh is made by folks like me,
then bridges find a spanning tree.

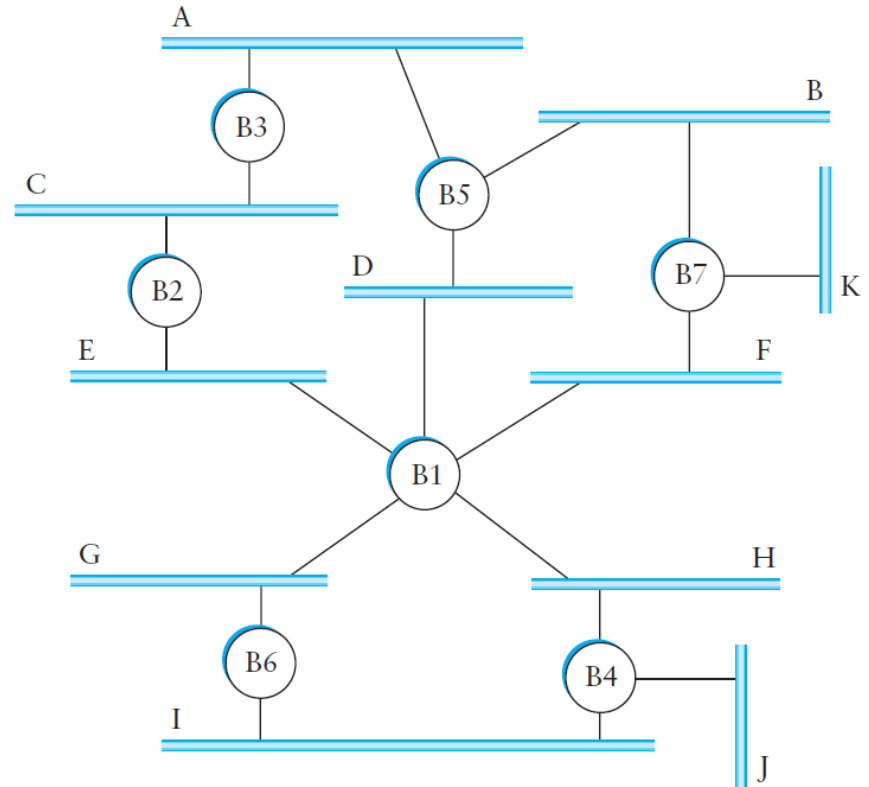
Radia Perlman



Spanning Tree Concepts

Key concepts:

- A single **root** bridge is elected
 - Each subnet must have a **single path** to reach the root bridge
- Each bridge may be connected to (and receive packets from) **multiple subnets**
 - Only the **designated bridge** will forward packets toward the root
- Every bridge knows which of its ports is **closest** to the root bridge
 - Called the **root port**



Conclusion

- Virtual Circuit Routing
- 802.11 Wireless
- Bridges and Spanning Tree Algorithm